Parameterized Approximation Schemes for Steiner Trees with Small Number of Steiner Vertices

Pavel Dvořák, Andreas Emil Feldmann, Dušan Knop,
Tomáš Masařík, Tomáš Toufar, Pavel Veselý

Faculty of Mathematics and Physics,
Charles University,
Prague, Czech Republic.

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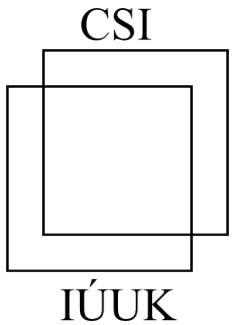
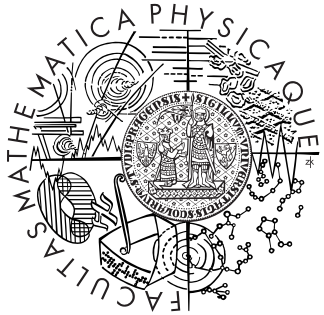
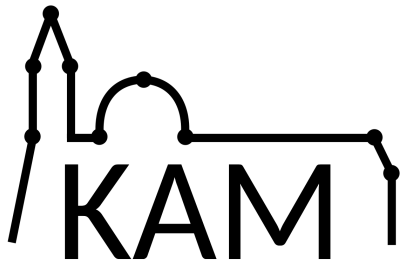


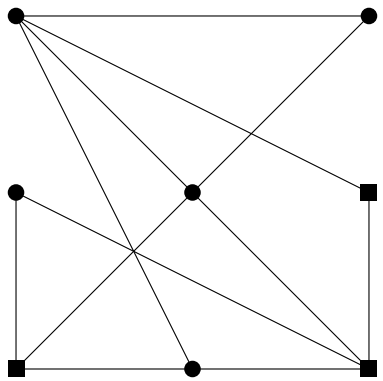




Steiner Tree

Steiner Tree
Input: A graph $G = (V, E)$, a set of terminals $R \subseteq V$.
Task: Find a tree $T \subseteq E$ of smallest size such that (R, T) is connected.

A vertex in $V \setminus R$ is called a Steiner vertex.
The number of Steiner vertices in the optimal solution is denoted by p, (i.e. $|V(T^*) \setminus R|$ for the optimum T^*).



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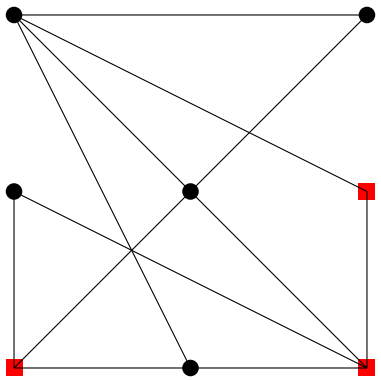

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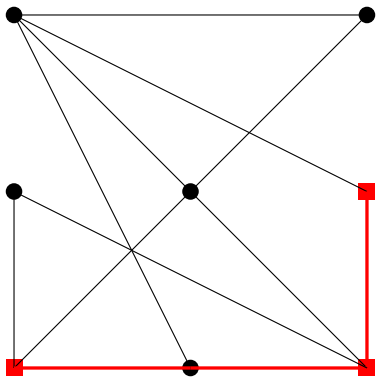

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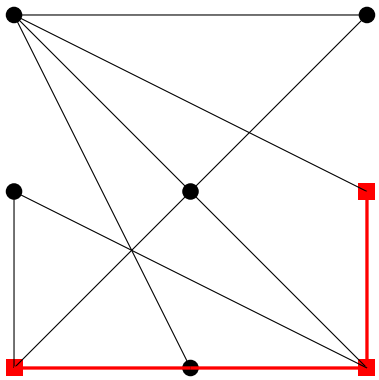

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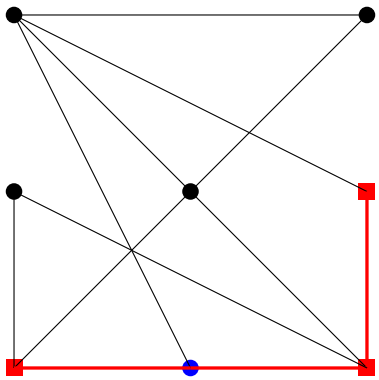

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Weighted Steiner Tree

Input: A graph $G = (V, E)$, a set of terminals $R \subseteq V$, and a weight function $w\colon E \to \mathbb{R}^+$.

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Directed Steiner Tree
Input: A directed graph $G = (V, A)$, a set of terminals $R \subseteq V$, a root $r \in R$.
Task: Find an arborescence $T \subseteq A$ of smallest size such that every terminal is reachable from r by edges in T.

Steiner Tree and Parameters

Two natural parameters of the problem.

- the number of terminals (denoted by $|R|$),
- the number of Steiner vertices in the optimal solution (denoted by p) i.e. $|V(T^*) \setminus R|$.

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Note, that if the number of Steiner vertices is considered as a parameter instead there is a trivial FPT algorithm.

Steiner Tree – Known Results

Parameterized complexity:

- FPT w.r.t. number of terminals
  - $3^k n^3$ algorithm in the 70s [Dreyfus & Wagner 71],
  - the best known now is $(2+\delta_\varepsilon)^k n^{o(1)}$ for any $\delta_\varepsilon \geq 0$ [Fuchs et al. 07] and $2^k n^2$ [Björklund et al. 07] for the unweighted case.
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Approximation:

- 2-approximation is simple
- $(\ln 4 + \varepsilon)$-approximation algorithm [Byrka et al. 13]
- (96/95)-approximation is NP-hard (which means no approximation scheme exists unless P = NP) [Chlebík and Chlebíková 02]

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Since Steiner Tree is both hard to approximate and W[1]-hard w.r.t. p, it is an excellent target for parameterized approximation.

Steiner Tree – Our Results

Theorem

For any computable functions g, f, it is impossible to compute an $f(p)$-approximation for the Weighted Directed Steiner Tree problem parameterized by p in time $g(p) \cdot n^{\mathrm{O}(1)}$, unless W[1] = FPT.

Results	Undirected	Directed
Unweighted		
Weighted		$f(p)$-approx. hard

Steiner Tree – Our Results

Definition (EPAS)

Efficient Parameterized Approximation Scheme for minimization problem is defined as follows: For all $\varepsilon > 0$ there exists an algorithm that runs in time $f(\varepsilon, p) \cdot n^{\mathrm{O}(1)}$ and returns $(1+\varepsilon)$-approximation.

Steiner Tree – Our Results

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There is an EPAS for the Unweighted Directed Steiner Tree problem

$$\text{in time } 2^{p^2/\varepsilon} \cdot n^{\mathrm{O}(1)}.$$

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Steiner Tree – Our Results

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- If not then we have a bounded number of terminals in terms of ε and p: $|R| \leq \frac{2p}{\varepsilon}$.
- So, we can use a known algorithm parameterized by $|R|$.

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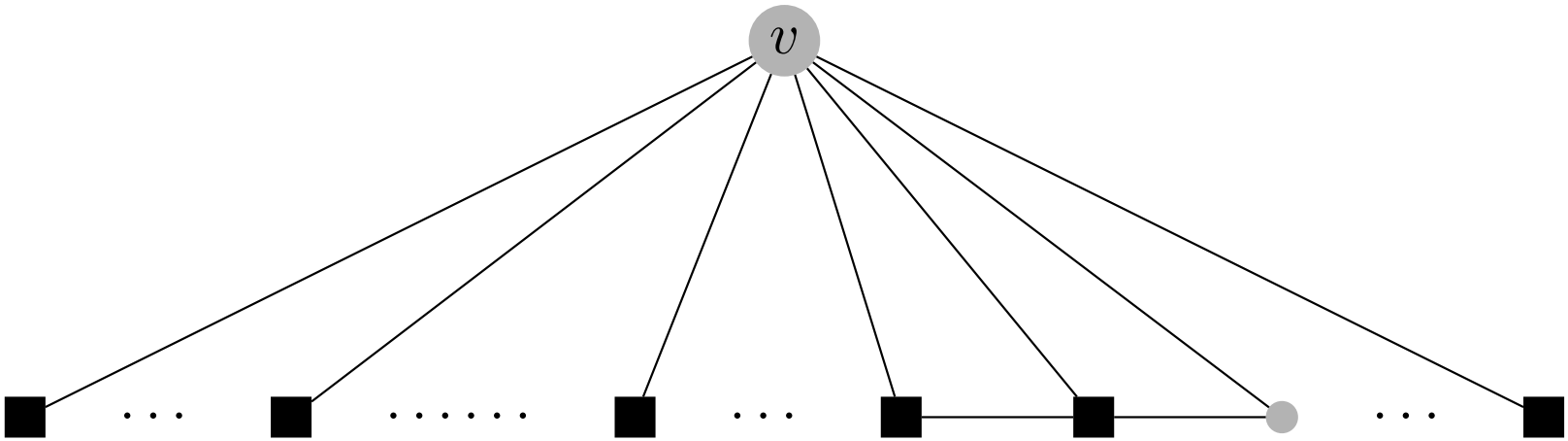
Unweighted		EPAS
Weighted	EPAS	$f(p)$-approx. hard

Overview of algorithm

We define a star C as the central vertex v and the subset Q of its adjacent terminals, where $|Q| \geq 2$.

We define a ratio of a star C, as

$$\frac{w(C)}{(|Q|-1)}.$$



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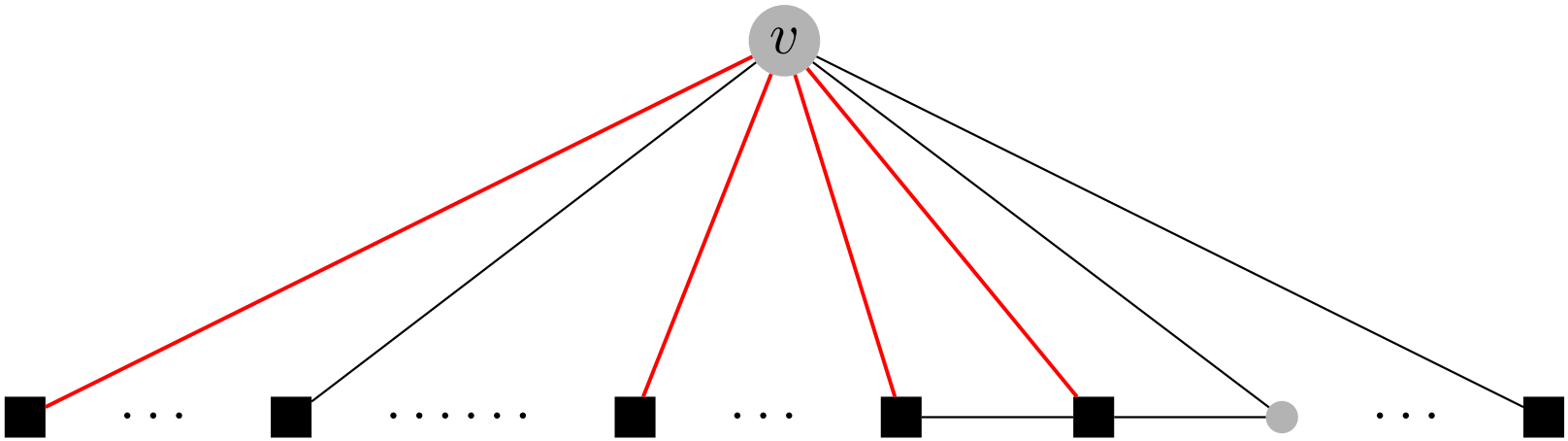

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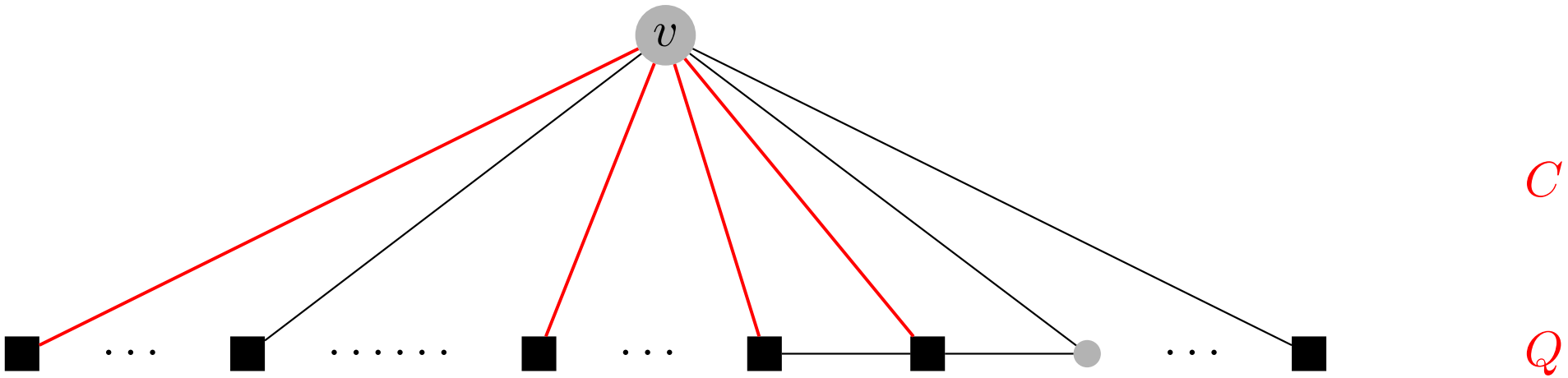

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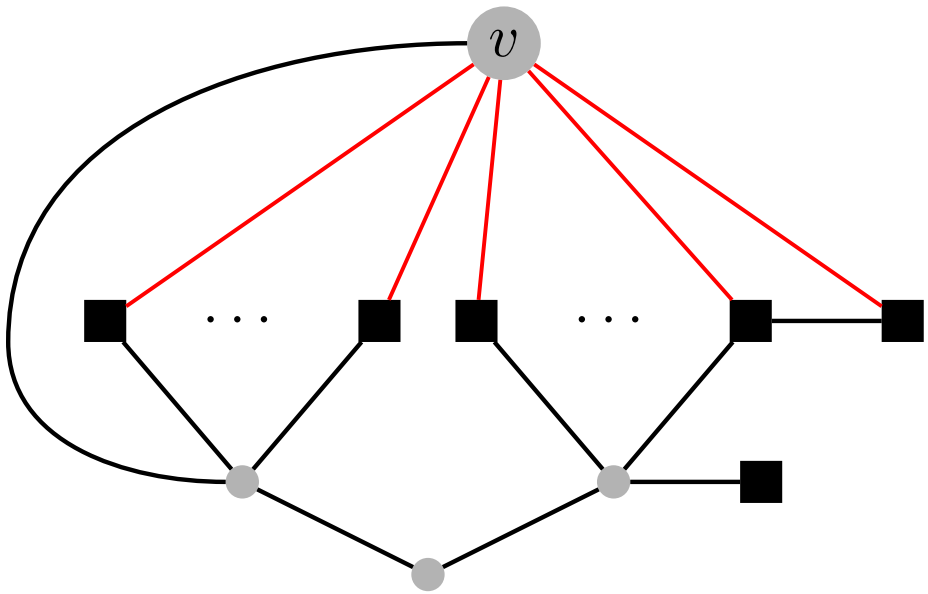

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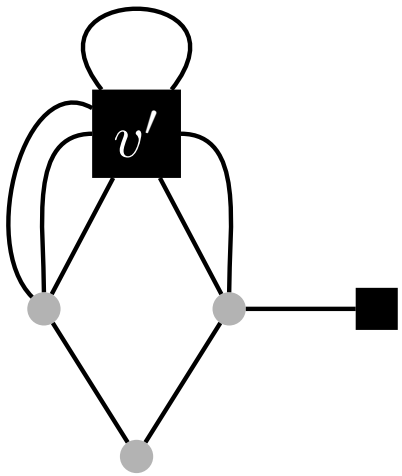

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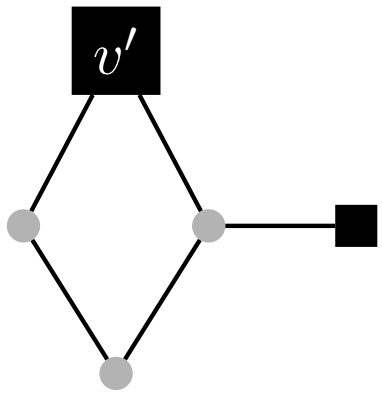

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- For analysis we start with an optimal solution T^*.
- We maintain a tree originating from T^* in the contracted instances.
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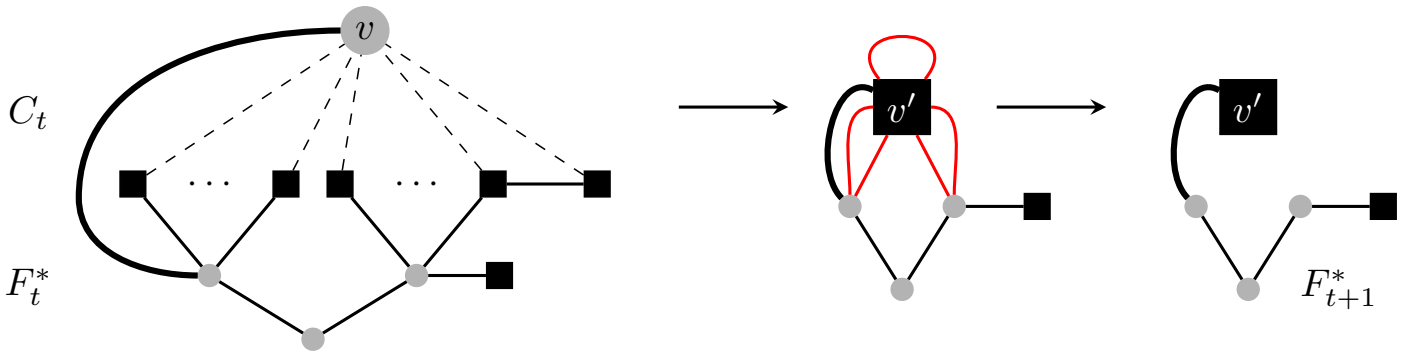

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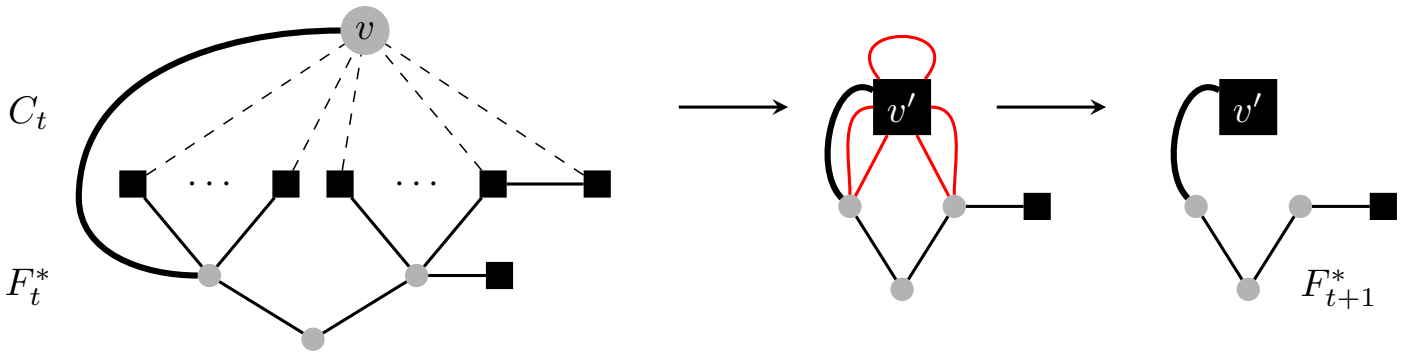
- In each step we compare weight of $w(C_t)$ with $w(D_t)$.
- D_t is a feedback edge set of T^*_t / C_t



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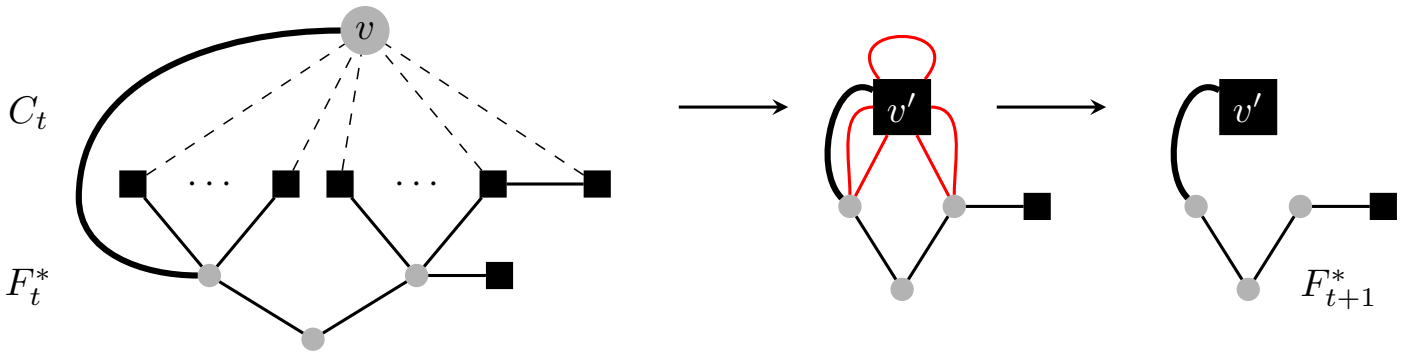

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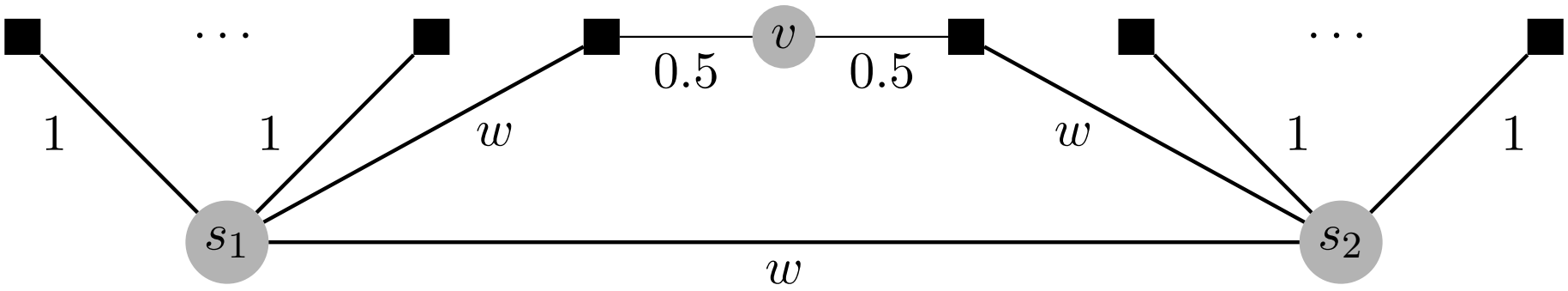

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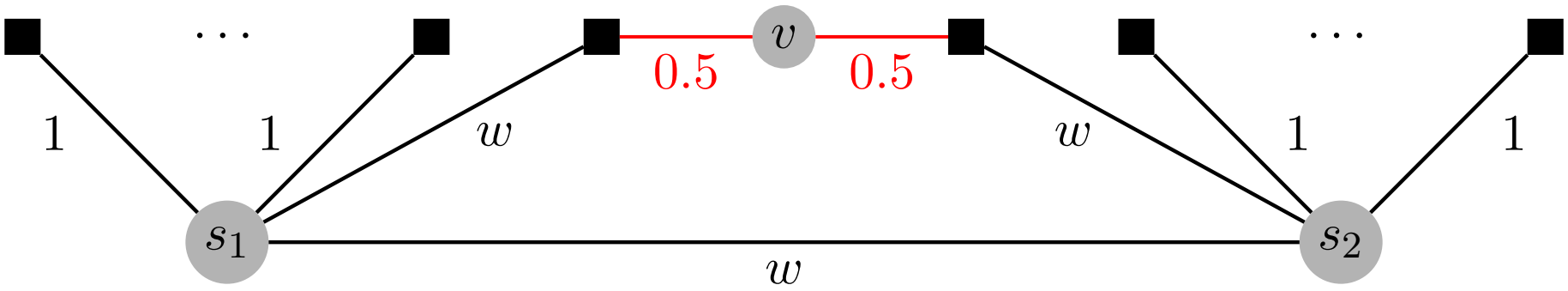

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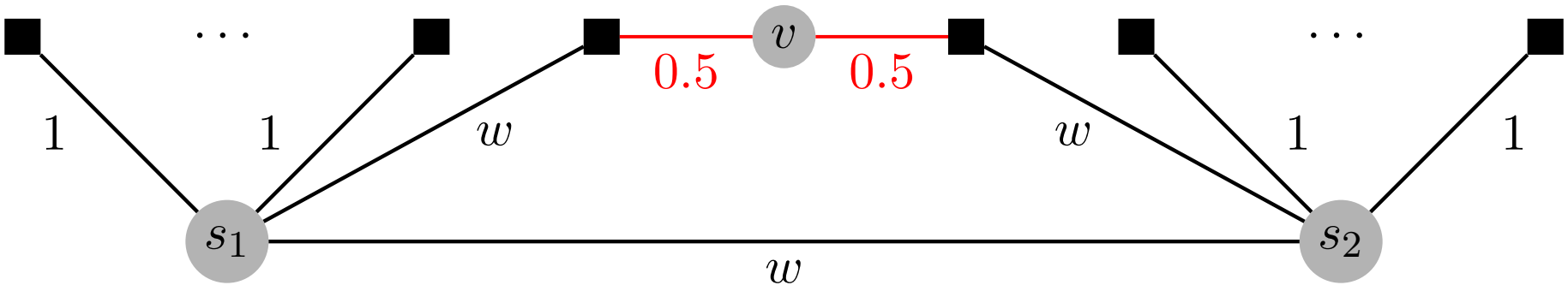

Proposition

If a star C contains at least $(1+\varepsilon)p/\varepsilon$ terminals then contraction of C is good.

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- We rescale weights to be ≥ 1.

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- otherwise to prove that the number of terminals is bounded in therms of p and ε.

Kernelization – Our Results

Definition (PSAKS)

Polynomial Size Approximate Kernelization Scheme for minimization problem is defined as follows: For all $\varepsilon > 0$ there exists an algorithm that runs in polynomial time ($n^{\mathrm{O}(1)}$) s.t. it returns an $(1+\varepsilon)$-approximate kernel of size $\mathrm{O}(p^{f(\varepsilon)})$.

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Theorem

No polynomial size $(2-\varepsilon)$-approximate kernelization algorithm exists for Unweighted Directed Steiner Tree parameterized by p for any $\varepsilon > 0$, unless NP $\subseteq$ coNP/poly.

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There is a PSAKS for the Weighted Undirected Steiner Tree problem. It computes an $(1+\varepsilon)$-approximate kernel

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- Combining these two we get a polynomial size $(1+\varepsilon)$-approximate kernel

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Generalizations of Steiner Tree

Weighted Steiner Forest
Input: A graph $G = (V, E)$, a list of terminal pairs $\{s_1, t_1\}, \ldots, \{s_k, t_k\}$, and a weight function $w\colon E \to \mathbb{R}^+$.
Task: Find a forest $F \subseteq E$ of smallest weight such that s_i and t_i is in the same component of F for every i.

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Runtime lower-bound

What is the best runtime dependence on p and ε under Exponential Time Hypothesis?

Open problems

EPSAKS	Undirected	Directed
Unweighted	unknown	no PSAKS
Weighted	unknown	$f(p)$-approx. hard

Runtime lower-bound

What is the best runtime dependence on p and ε under Exponential Time Hypothesis?

Thank you for your attention!